

gamete. If $t[j] = 0$, then there exists some taxon t' for which $t'[j] = 0$ and $t'[i] = 1$ and therefore $t'[i'] = 1$ and again $(0, 1)$ cannot be a newly introduced gamete. Finally consider any pair of characters j, j' . If taxon t introduces gamete $(0, 1)$, then there exists some taxon t' with $t'[j] = 0$ and $t'[i] = 1$. If $t'[j'] = 1$, then $(0, 1)$ cannot be a new gamete. If $t'[j'] = 0$, then $t[j'] = 0$ and not 1. The case when $(1, 0)$ is introduced by t is symmetric. Finally if t introduces $(1, 1)$ then consider any taxon t' with $t'[i] = 1$. It has to be the case that $t'[j] = t'[j'] = 1$, and therefore $(1, 1)$ cannot be a newly introduced gamete. ■

We now have the following lemma.

Lemma 4.5: In every optimum phylogeny T_M^* , the conflict graph on the set of taxa in T_M^* (Steiner vertices included) is the same as the conflict graph on M .

Proof: We say that a subgraph F' of F is the same as an edge labeled tree T if F' is a tree and T can be obtained from F' by suppressing degree-two vertices. A phylogeny T is contained in a graph F if there exists an edge-labeled subgraph F' that is the same as the edge labeled (by function μ) phylogeny T . We know from Lemma 4.3 that all optimum phylogenies T_M^* for M is contained in the (extended) Buneman graph of M . Lemma 4.4 shows that the conflict graph on M'' (and therefore on the extended Buneman graph of M'') is the same as the conflict graph of M . ■

Lemma 4.6: The probability that all guesses performed at Step 2c are correct is at least 2^{-q} .

Proof: Implementation: We first show how to perform the guess efficiently. For every character i , we perform the following steps in order.

- 1) if all taxa in M_0 contain the same state s in i , then fix $r[i] = s$
- 2) if all taxa in M_1 contain the same state s in i , then fix $r[i] = s$
- 3) if $r[i]$ is unfixed then guess $r[i]$ uniformly at random from $\{0, 1\}$

Assuming that the guess at Step 2a (Figure 4) is correct, we know that there exists an optimum phylogeny $T_{M_j}^*$ on M_j where $c(v)$ mutates exactly once. Let $e \in T_{M_j}^*$ s.t. $c(v) \in \mu(e)$. Let r' be an end point of e s.t. $r'[c(v)] = 1$ and p' be the other end point. If the first two conditions hold with the same state s , then character i does not mutate in M_j . In such a case, we know that $r'[i] = s$, since $T_{M_j}^*$ is optimal and the above method ensures that $r[i] = s$. Notice that if both conditions are satisfied simultaneously with different values of s then i and $c(v)$ share exactly two gametes in M_j and therefore $i, c(v) \in \mu(e)$. Hence, $r'[i] = r[i]$. We now consider the remaining cases when exactly one of the above conditions hold. We show that if $r[i]$ is fixed